

## SQL Recursion

WITH

stuff that looks like Datalog rules  
an SQL query about EDB, IDB

- Rule =

[RECURSIVE] *R*(<arguments>) AS  
SQL query

## Example

Find Sally's cousins, using EDB Par(child, parent).

```
WITH
  Sib(x,y) AS
    SELECT p1.child, p2.child
    FROM Par p1, Par p2
    WHERE p1.parent = p2.parent
          AND p1.child <> p2.child,
  RECURSIVE Cousin(x,y) AS
    Sib
    UNION
    (SELECT p1.child, p2.child
    FROM Par p1, Par p2, Cousin
    WHERE p1.parent = Cousin.x
          AND p2.parent = Cousin.y
    )

SELECT y
FROM Cousin
WHERE x = 'Sally';
```

## Plan for Describing Legal SQL recursion

1. Define “monotonicity,” a property that generalizes “stratification.”
2. Generalize stratum graph to apply to SQL queries instead of Datalog rules.
  - ❖ (Non)monotonicity replaces NOT in subgoals.
3. Define semantically correct SQL recursions in terms of stratum graph.

### Monotonicity

If relation  $P$  is a function of relation  $Q$  (and perhaps other things), we say  $P$  is *monotone* in  $Q$  if adding tuples to  $Q$  cannot cause any tuple of  $P$  to be deleted.

## Monotonicity Example

In addition to certain negations, an aggregation can cause nonmonotonicity.

```
Sells(bar, beer, price)
```

```
SELECT AVG(price)  
FROM Sells  
WHERE bar = 'Joe's Bar';
```

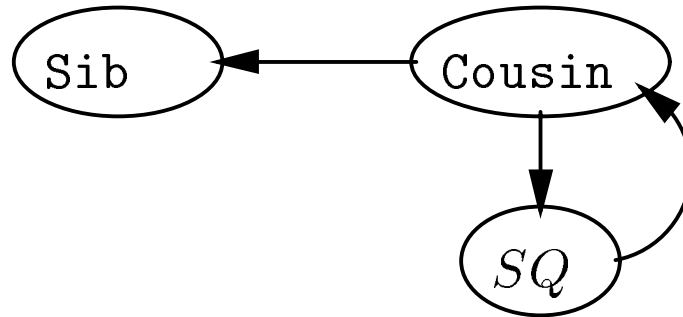
- Adding to `Sells` a tuple that gives a new beer Joe sells will usually change the average price of beer at Joe's.
- Thus, the former result, which might be a single tuple like (2.78) becomes another single tuple like (2.81), and the old tuple is lost.

## Generalizing Stratum Graph to SQL

- Node for each relation defined by a “rule.”
- Node for each subquery in the “body” of a rule.
- Arc  $P \rightarrow Q$  if
  - a)  $P$  is “head” of a rule, and  $Q$  is a relation appearing in the FROM list of the rule (not in the FROM list of a subquery), as argument of a UNION, etc.
  - b)  $P$  is head of a rule, and  $Q$  is a subquery directly used in that rule (not nested within some larger subquery).
  - c)  $P$  is a subquery, and  $Q$  is a relation or subquery used directly within  $P$  [analogous to (a) and (b) for rule heads].
- Label the arc – if  $P$  is *not* monotone in  $Q$ .
- Requirement for legal SQL recursion: finite strata only.

## Example

For the Sib/Cousin example, there are three nodes: Sib, Cousin, and  $SQ$  (the second term of the union in the rule for Cousin).



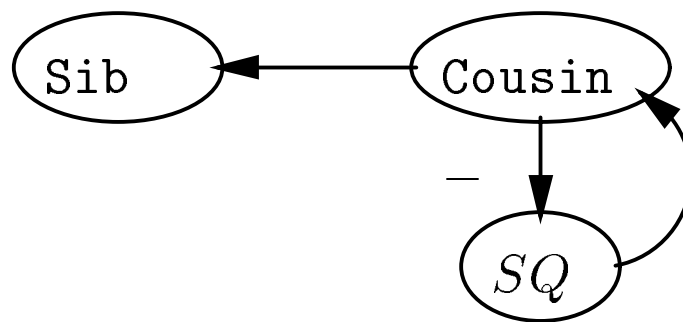
- No nonmonotonicity, hence legal.

## A Nonmonotonic Example

Change the UNION to EXCEPT in the rule for Cousin.

```
RECURSIVE Cousin(x,y) AS
  Sib
  EXCEPT
  (SELECT p1.child, p2.child
   FROM Par p1, Par p2, Cousin
   WHERE p1.parent = Cousin.x
        AND p2.parent = Cousin.y
  )
```

- Now, adding to the result of the subquery can delete Cousin facts; i.e., Cousin is nonmonotone in  $SQ$ .



- Infinite number of  $-$ 's in cycle, so illegal in SQL.

## Another Example: NOT Doesn't Mean Nonmonotone

Leave `Cousin` as it was, but negate one of the conditions in the where-clause.

```
RECURSIVE Cousin(x,y) AS
  Sib
  UNION
  (SELECT p1.child, p2.child
   FROM Par p1, Par p2, Cousin
   WHERE p1.parent = Cousin.x
        AND NOT (p2.parent = Cousin.y)
  )
```

- You might think that  $SQ$  depends negatively on `Cousin`, but it doesn't.
  - ❖ If I add a new tuple to `Cousin`, all the old tuples still exist and yield whatever tuples in  $SQ$  they used to yield.
  - ❖ In addition, the new `Cousin` tuple might combine with old  $p1$  and  $p2$  tuples to yield something new.



## Object-Oriented DBMS's

- ODMG = Object Data Management Group: an OO standard for databases.
- ODL = Object Description Language: design in the OO style.
- OQL = Object Query Language: queries an OO database with an ODL schema, in a manner similar to SQL.

## ODL Overview

Class declarations include:

1. Name for the class.
2. Key declaration(s), which are optional.
3. *Extent* declaration = name for the set of currently existing objects of a class.
4. *Element* declarations. An element is an attribute, a relationship, or a method.

## ODL Class Declarations

```
class <name> {  
    elements = attributes, relationships,  
             methods  
}
```

## Element Declarations

```
attribute <type> <name>;  
relationship <rangetype> <name>;
```

- Relationships involve objects; attributes (usually) involve non-object values, e.g., integers.

## Method Example

```
float gpa(in string) raises(noGrades)
```

- `float` = return type.
- `in`: indicates the argument (a student name, presumably) is read-only.
  - ❖ Other options: `out`, `inout`.
- `noGrades` is an exception that can be raised by method `gpa`.

## ODL Relationships

- Only binary relations supported.
  - ❖ Multiway relationships require a “connecting” class, as discussed for E/R model.
- Relationships come in inverse pairs.
  - ❖ Example: “Sells” between beers and bars is represented by a relationship in bars, giving the beers sold, *and* a relationship in beers giving the bars that sell it.
- Many-many relationships have a set type (called a *collection type*) in each direction.
- Many-one relationships have a set type for the one, and a simple class name for the many.
- One-one relations have classes for both.

## Beers-Bars-Drinkers Example

```
class Beers {
    attribute string name;
    attribute string manf;
    relationship Set<Bars> servedAt
        inverse Bars::serves;
    relationship Set<Drinkers> fans
        inverse Drinkers::likes;
}
```

- An element from another class is indicated by <class>::
- Form a set type with Set<type>.

```
class Bars {
    attribute string name;
    attribute Struct Addr
        {string street, string city, int zip}
        address;
    attribute Enum Lic {full, beer, none}
        licenseType;
    relationship Set<Drinkers> customers
        inverse Drinkers::frequents;
    relationship Set<Beers> serves
        inverse Beers::servedAt;
}
```

- Structured types have names and bracketed lists of field-type pairs.
- Enumerated types have names and bracketed lists of values.

```
class Drinkers {  
    attribute string name;  
    attribute Struct Bars::Addr  
        address;  
    relationship Set<Beers> likes  
        inverse Beers::fans;  
    relationship Set<Bars> frequents  
        inverse Bars::customers;  
}
```

- Note reuse of Addr type.

## ODL Type System

- Basic types: int, real/float, string, enumerated types, and classes.
- Type constructors: **Struct** for structures and five *collection types*: **Set**, **Bag**, **List**, **Array**, and **Dictionary**.
- Relationship types may only be classes or a collection of a class.



## Many-One Relationships

Don't use a collection type for relationship in the "many" class.

### Example: Drinkers Have Favorite Beers

```
class Drinkers {
    attribute string name;
    attribute Struct Bars::Addr
        address;
    relationship Set<Beers> likes
        inverse Beers::fans;
    relationship Beers favoriteBeer
        inverse Beers::realFans;
    relationship Set<Bars> frequents
        inverse Bars::customers;
}
```

- Also add to Beers:

```
relationship Set<Drinkers> realFans
    inverse Drinkers::favoriteBeer;
```

## Example: Multiway Relationship

Consider a 3-way relationship bars-beers-prices. We have to create a connecting class BBP.

```
class Prices {
    attribute real price;
    relationship Set<BBP> toBBP
        inverse BBP::thePrice;
}

class BBP {
    relationship Bars theBar inverse ...
    relationship Beers theBeer inverse ...
    relationship Prices thePrice
        inverse Prices::toBBP;
}
```

- Inverses for `theBar`, `theBeer` must be added to `Bars`, `Beers`.
- Better in this special case: make no `Prices` class; make `price` an attribute of `BBP`.
- Notice that keys are optional.
  - ❖ `BBP` has no key, yet is not “weak.” Object identity suffices to distinguish different `BBP` objects.

## Roles in ODL

Names of relationships handle “roles.”

### Example: Spouses and Drinking Buddies

```
class Drinkers {
    attribute string name;
    attribute Struct Bars::Addr
        address;
    relationship Set<Beers> likes
        inverse Beers::fans;
    relationship Set<Bars> frequents
        inverse Bars::customers;
    relationship Drinkers husband
        inverse wife;
    relationship Drinkers wife
        inverse husband;
    relationship Set<Drinkers> buddies
        inverse buddies;
}
```

- Notice that `Drinkers::` is optional when the inverse is a relationship of the same class.