



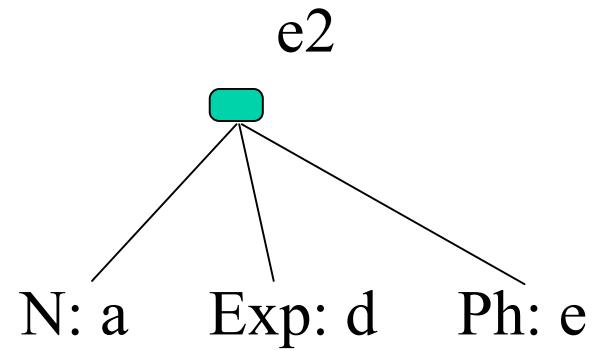
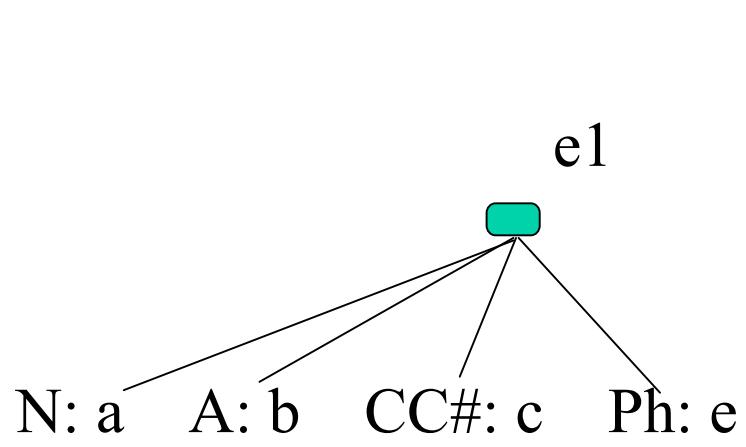
Evaluating Entity Resolution Results

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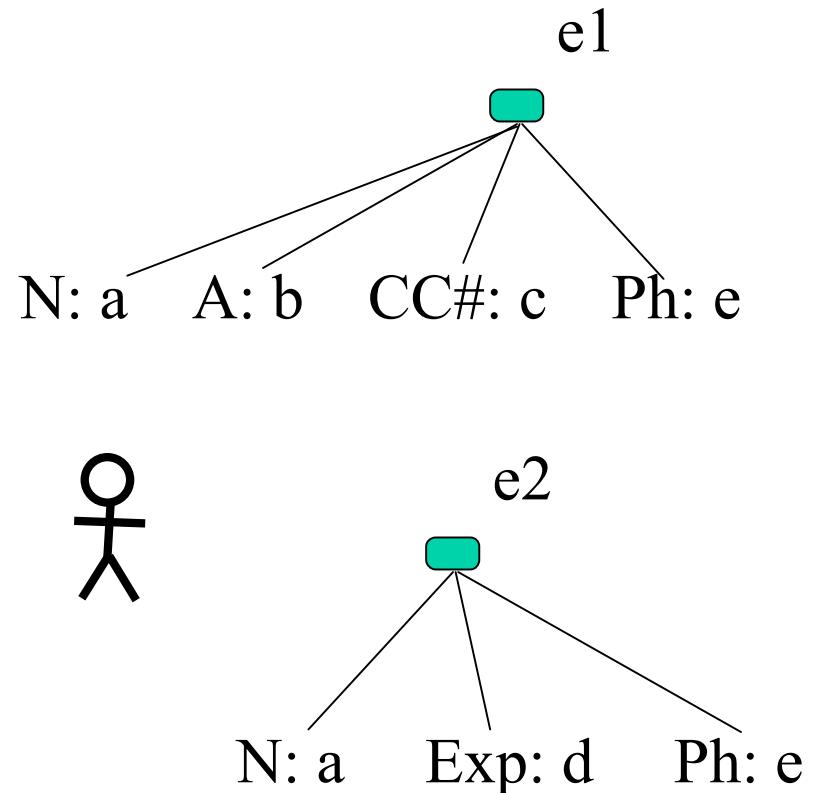
Stanford University

Entity Resolution



Applications

- comparison shopping
- mailing lists
- classified ads
- customer files
- counter-terrorism



Evaluating ER Results

R1 = a, b, c, d, efgh

R2 = ab, cd, ef, gh

G = ab, cd, efgh

Pairwise Recall

R1 = a, b, c, d, e f g h

R2 = ab, cd, ef, gh

G = ab, cd, e f g h

Pairwise Recall

R1 = a, b, c, d, e f g h

Pairs:

ef, eg, eh,
fg, fh, gh

R2 = ab, cd, ef, gh

Pairs:

ab, cd, ef, gh

G = ab, cd, e f g h

Pairs:

ab, cd, ef, eg,
eh, fh, fh, gh

Pairwise Recall

R1 = a, b, c, d, e f g h

R2 = ab, cd, ef, gh

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Pairs:

ef, eg, eh,
fg, fh, gh

Pairs:

ab, cd, ef, gh

Pairs:

ab, cd, ef, eg,
eh, fh, fh, gh

6 pairs, all in G

8 pairs

Pairwise Recall

R1 = a, b, c, d, e f g h

R2 = ab, cd, ef, gh

G = ab, cd, e f g h

Pairs:

ef, eg, eh,
fg, fh, gh

Pairs:

ab, cd, ef, gh

Pairs:

ab, cd, ef, eg,
eh, fh, fh, gh

6 pairs, all in G

8 pairs

Recall = 6/8 = 75%

Pairwise Recall

$R_1 = a, b, c, d, e, f, g, h$

Pairs:

ef, eg, eh,
fg, fh, gh

6 pairs, all in G

Recall = $6/8 = 75\%$

$R_2 = a, b, c, d, e, f, g, h$

Pairs:

ab, cd, ef, gh

4 pairs, all in G

Recall = $4/8 = 50\%$

$G = a, b, c, d, e, f, g, h$

Pairs:

ab, cd, ef, eg,
eh, fh, fh, gh

8 pairs

Pairwise Recall

R1 = a, b, c, d, efg

R2 = ab, cd, ef, gh

G = ab, cd, efg

Pairs:

ef, eg, eh,
fg, fh, gh

Pairs:

ab, cd, ef, gh

Pairs:

ab, cd, ef, eg,
eh, fh, fh, gh

6 pairs, all in G

4 pairs, all in G

8 pairs

Recall = $6/8 = 75\%$

Recall = $4/8 = 50\%$

Pairwise F1

$$PairPrecision(R, G) = \frac{|Pairs(R) \cap Pairs(G)|}{|Pairs(R)|}$$

$$PairRecall(R, G) = \frac{|Pairs(R) \cap Pairs(G)|}{|Pairs(G)|}$$

$$pF_1 = \frac{2 \times Precision \times Recall}{Precision + Recall}$$

Merge Distance

R1 = a, b, c, d, efgh

R2 = ab, cd, ef, gh

G = ab, cd, efgh

Merge Distance

R1 = a, b, c, d, e f g h

R2 = ab, cd, ef, gh

G = ab, cd, e f g h

a, b → ab

c, d → cd

G = ab, cd, e f g h

Merge Distance

R1 = a, b, c, d, e f g h

R2 = a b, c d, e f, g h

G = a b, c d, e f g h

a, b → ab
c, d → cd

e f, g h → e f g h

G = a b, c d, e f g h

G = a b, c d, e f g h

Merge Distance

R1 = a, b, c, d, efg

R2 = ab, cd, ef, gh

G = ab, cd, efg

a, b → ab
c, d → cd

ef, gh → efg

G = ab, cd, efg

G = ab, cd, efg

Distance = 2

Distance = 1

Merge Distance

Minimum number of splits and merges to
get from R to G (splits first)

[Al-Kamha, et al. 2004]

Variation of Information

$$VI(R, G) = H(R) + H(G) - 2I(R, G)$$

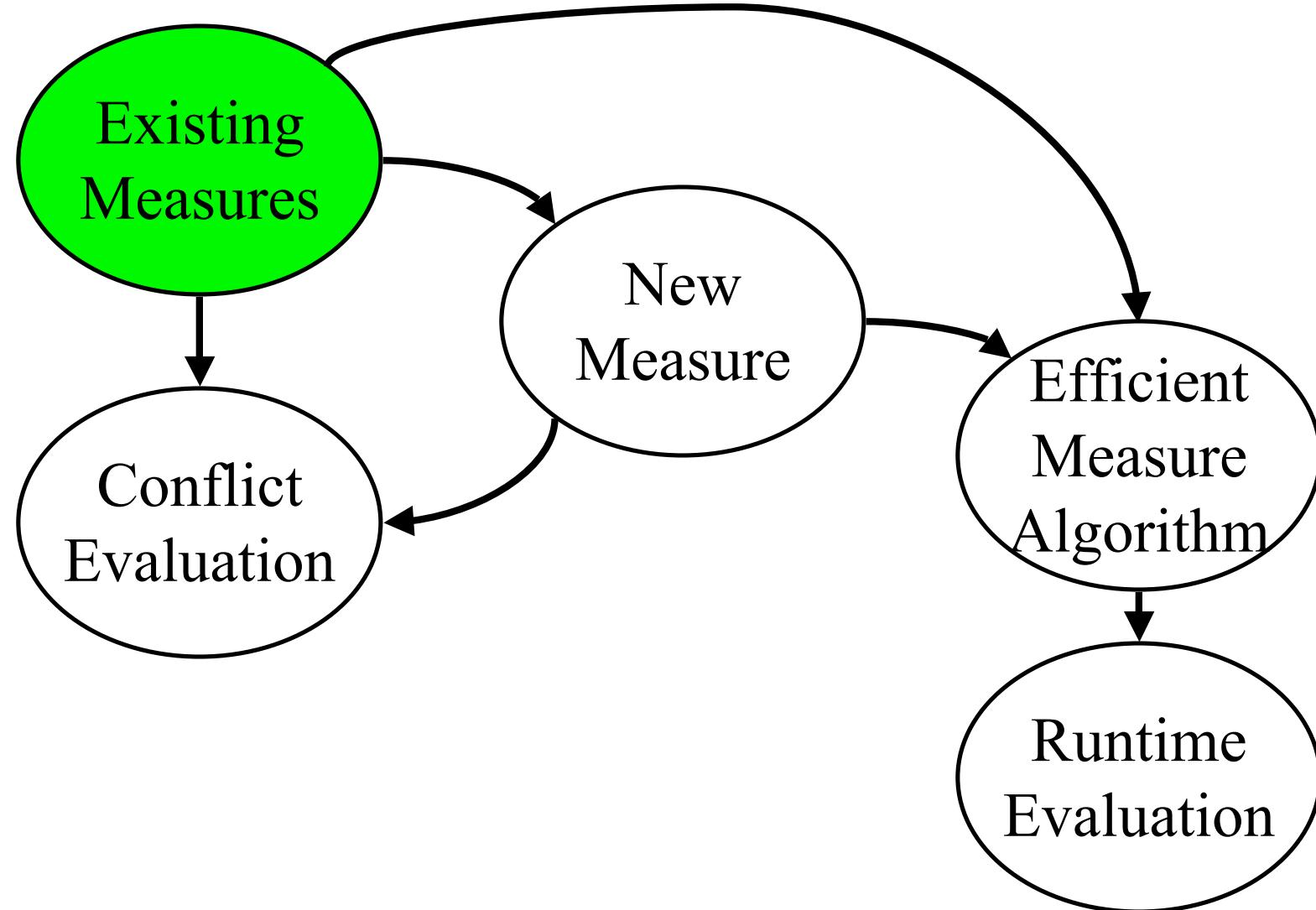
$$H(R) = - \sum_{r \in R} \frac{|r|}{N} \log \frac{|r|}{N}$$

$$I(R, G) = \sum_{r \in R} \sum_{g \in G} \frac{|r \cap g|}{N} \log \frac{|r \cap g| \times N}{|r| \times |g|}$$

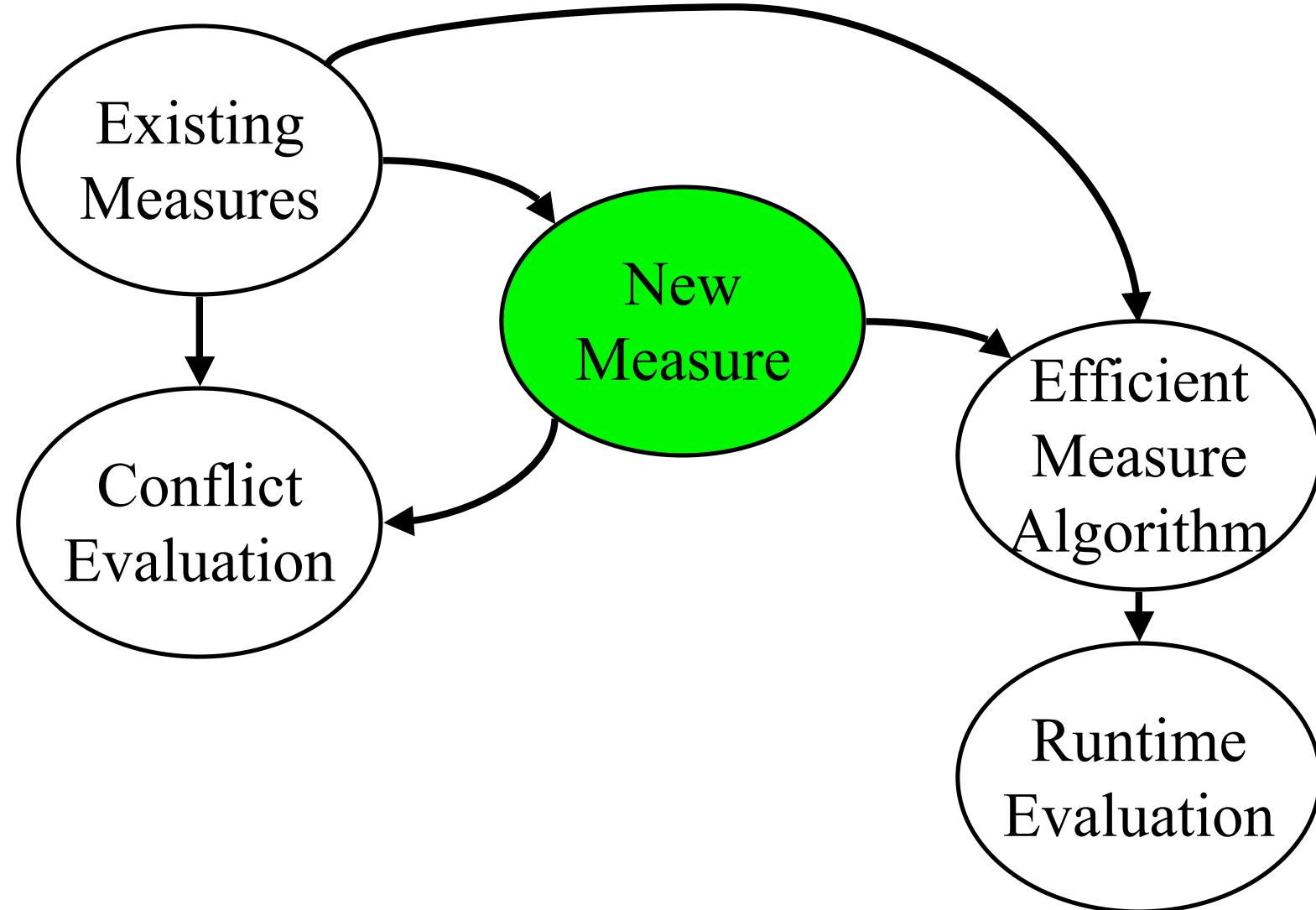
Conflicts

	R1	R2
Pairwise Recall	75%	50%
Merge Distance	2	1
Variation of Information	0.5	0.5

Road Map



Road Map



Generalized Merge Distance

- Cost of split, merge defined by functions:

$$f_s(x, y), f_m(x, y)$$

e.g.,

$$f_m(x, y) = 1$$

$$f_m(x, y) = xy$$

- Distance = cost of minimum-cost path

Generalized Merge Distance

R1 = a, b, c, d, efg

R2 = ab, cd, ef, gh

G = ab, cd, efg

a, b → ab
c, d → cd

ef, gh → efg

G = ab, cd, efg

G = ab, cd, efg

Distance
 $= f_m(1, 1) + f_m(1, 1)$

Distance
 $= f_m(2, 2)$

$$\underline{f(x, y) = 1}$$

$$R1 = a, b, c, d, efg$$

$$R2 = ab, cd, ef, gh$$

$$G = ab, cd, efg$$

$$\begin{aligned} a, b &\rightarrow ab \\ c, d &\rightarrow cd \end{aligned}$$

$$ef, gh \rightarrow efg$$

$$G = ab, cd, efg$$

$$G = ab, cd, efg$$

$$\begin{aligned} \text{Distance} \\ &= f_m(1, 1) + f_m(1, 1) \\ &= 1 + 1 = 2 \end{aligned}$$

$$\begin{aligned} \text{Distance} \\ &= f_m(2, 2) \\ &= 1 \end{aligned}$$

$$\underline{f(x, y) = xy}$$

$$R1 = a, b, c, d, efg$$

$$R2 = ab, cd, ef, gh$$

$$G = ab, cd, efg$$

$$\begin{aligned} a, b &\rightarrow ab \\ c, d &\rightarrow cd \end{aligned}$$

$$ef, gh \rightarrow efg$$

$$G = ab, cd, efg$$

$$G = ab, cd, efg$$

$$\begin{aligned} \text{Distance} \\ = f_m(1, 1) + f_m(1, 1) \end{aligned}$$

$$\begin{aligned} \text{Distance} \\ = f_m(2, 2) \end{aligned}$$

$$\underline{f(x, y) = xy}$$

$$R1 = a, b, c, d, efg$$

$$R2 = ab, cd, ef, gh$$

$$G = ab, cd, efg$$

$$\begin{aligned} a, b &\rightarrow ab \\ c, d &\rightarrow cd \end{aligned}$$

$$ef, gh \rightarrow efg$$

$$G = ab, cd, efg$$

$$G = ab, cd, efg$$

$$\begin{aligned} \text{Distance} \\ &= f_m(1, 1) + f_m(1, 1) \\ &= 1 \times 1 + 1 \times 1 = 2 \end{aligned}$$

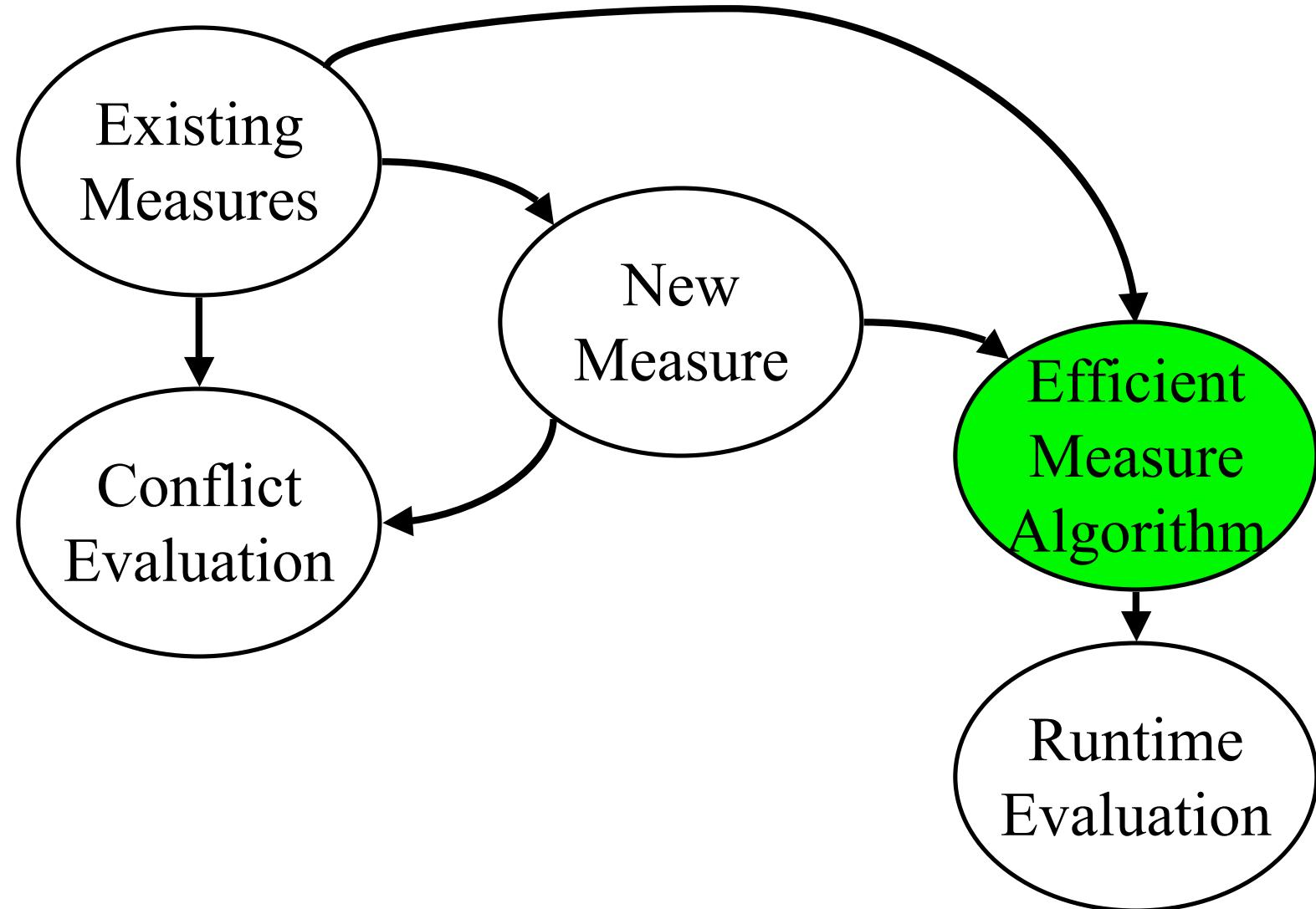
$$\begin{aligned} \text{Distance} \\ &= f_m(2, 2) \\ &= 2 \times 2 = 4 \end{aligned}$$

Relationships Between Measures

- Merge Distance: $f_m(x, y) = 1, f_s(x, y) = 1$
- Pairwise Recall: $f_m(x, y) = xy, f_s(x, y) = 0$
- Pairwise Precision: $f_m(x, y) = 0, f_s(x, y) = xy$
- Variation of Information:
$$f_m(x, y) = f_s(x, y) = h(x + y) - h(x) - h(y)$$

$$h(x) = \frac{x}{N} \log \frac{x}{N}$$

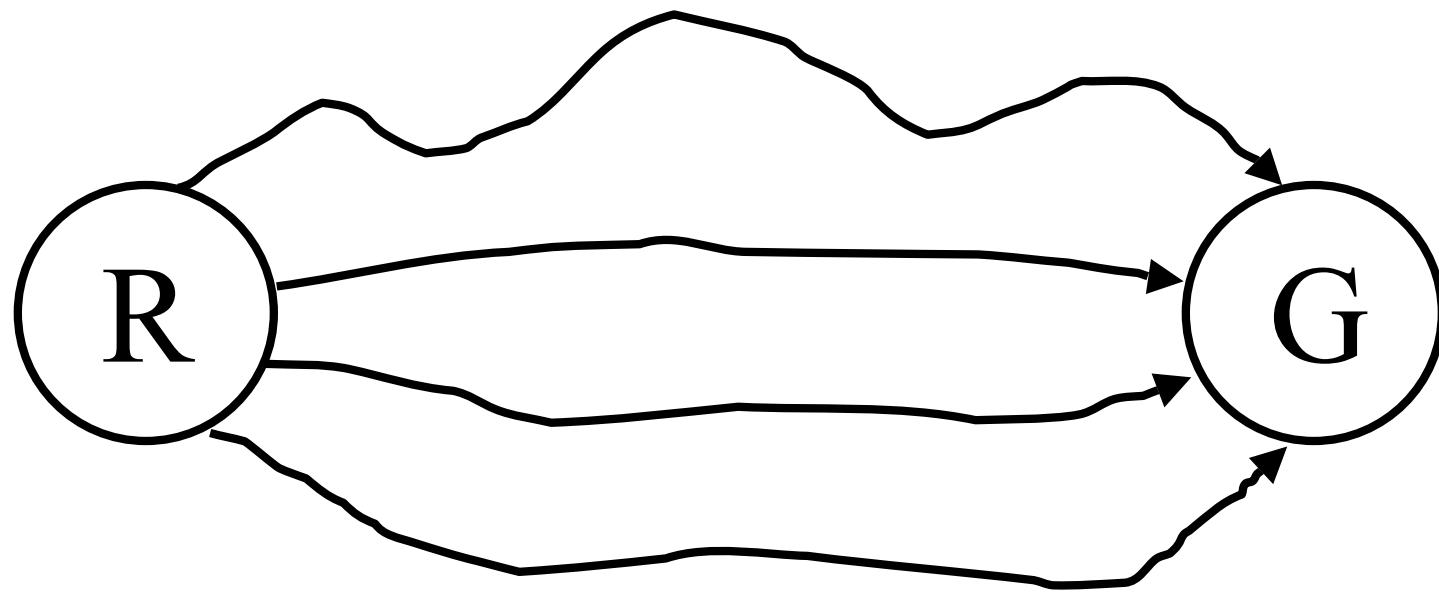
Road Map



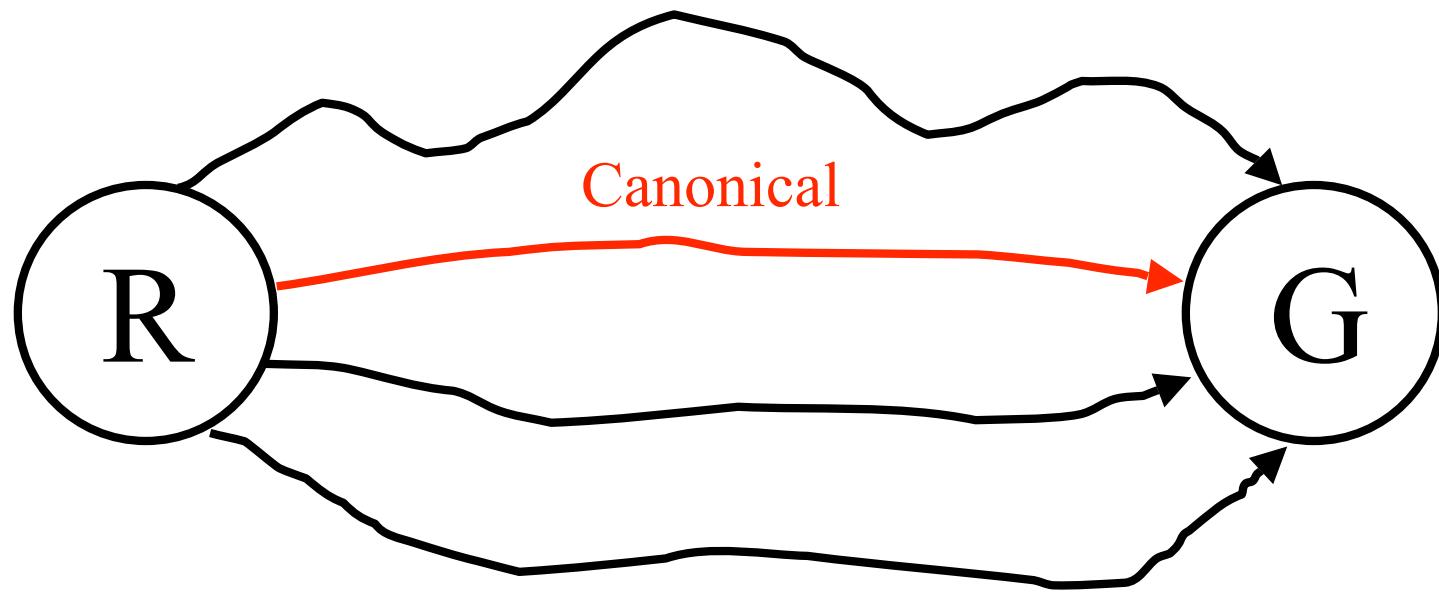
Slice Algorithm

- Linear time algorithm
- Extra property required:
$$f(x, y) + f(x+y, z) = f(x, z) + f(x+z, y)$$
- Cost functions for pairwise, merge distance, and variation of information (and many others) satisfy property

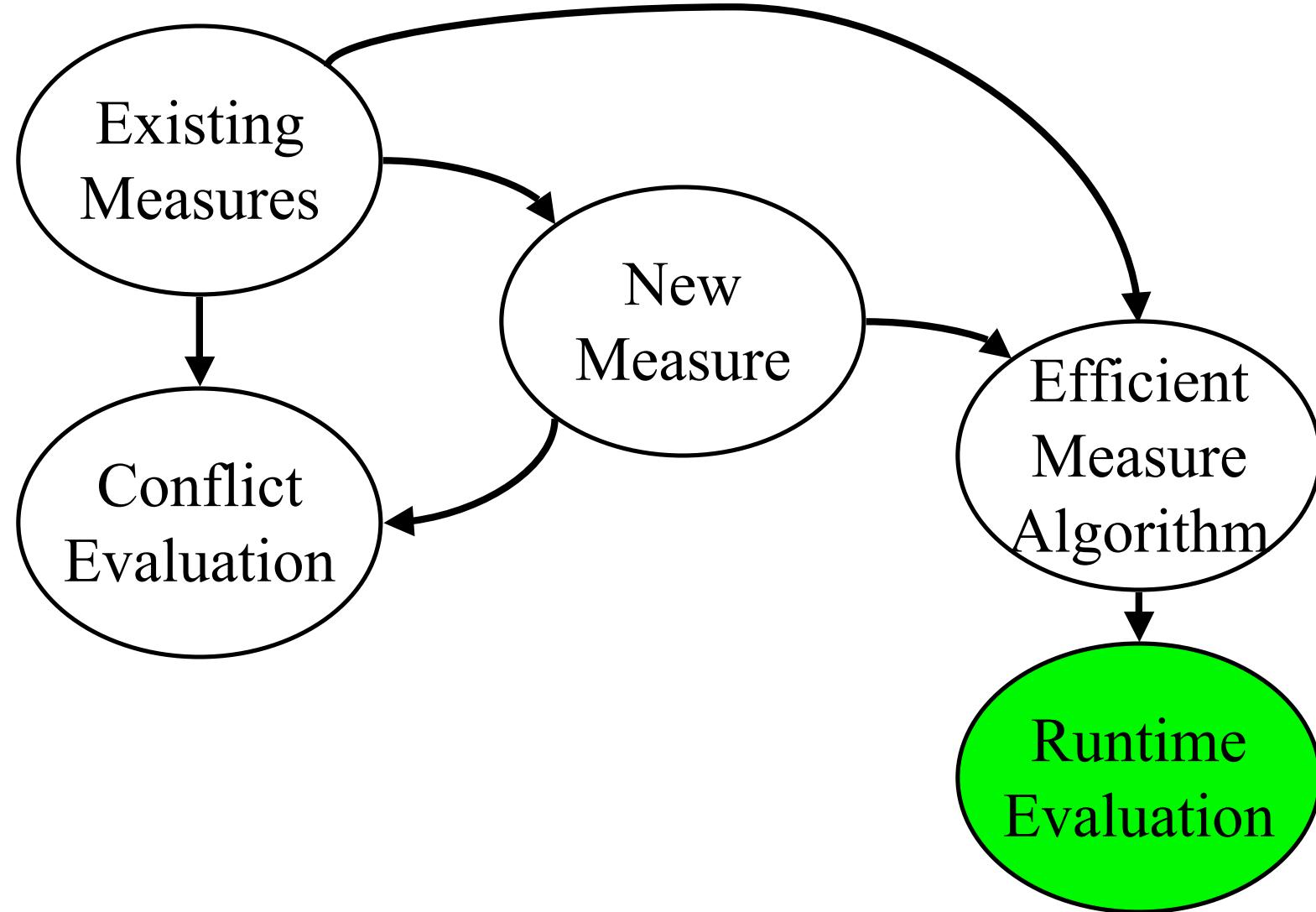
Slice Algorithm



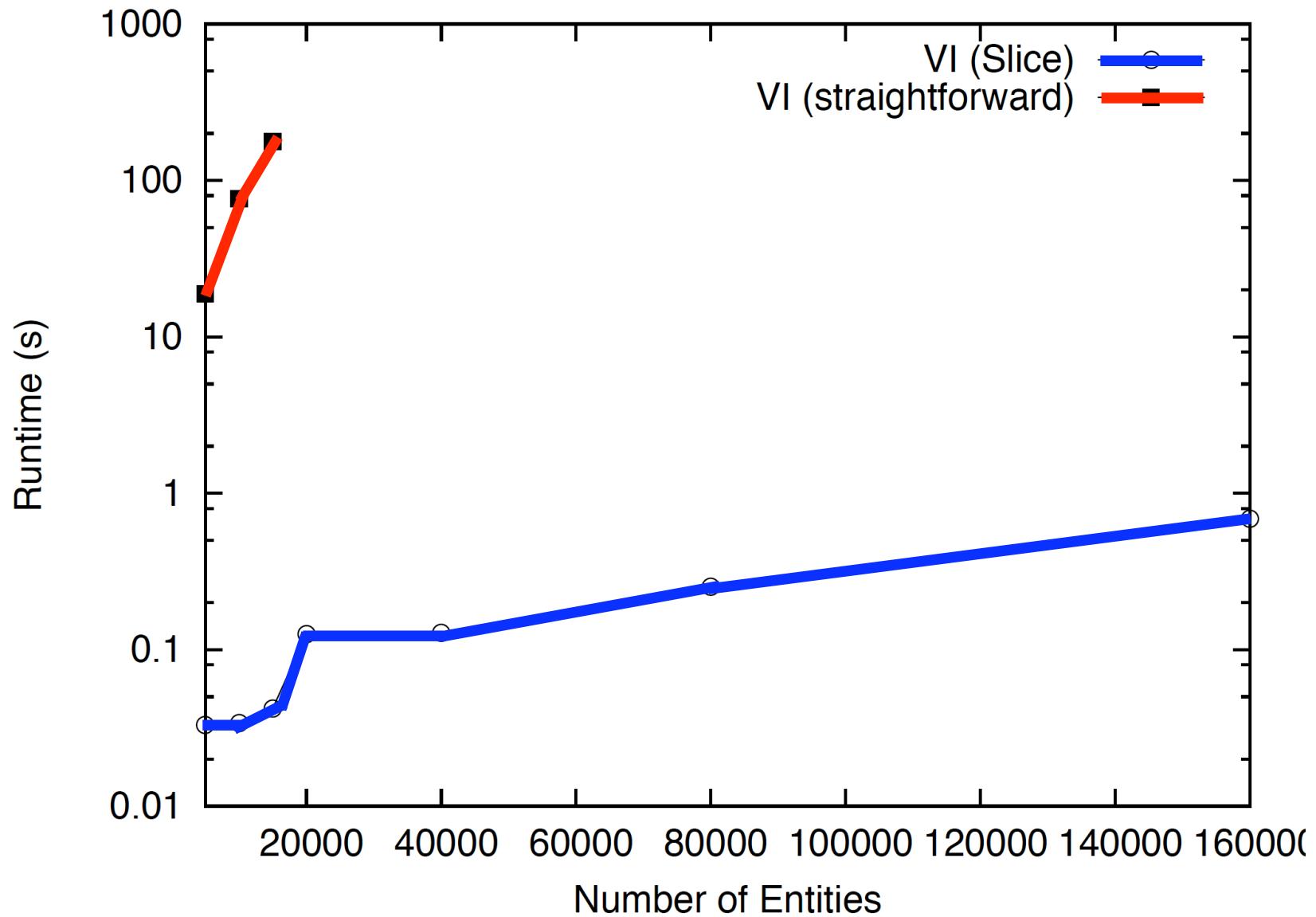
Slice Algorithm



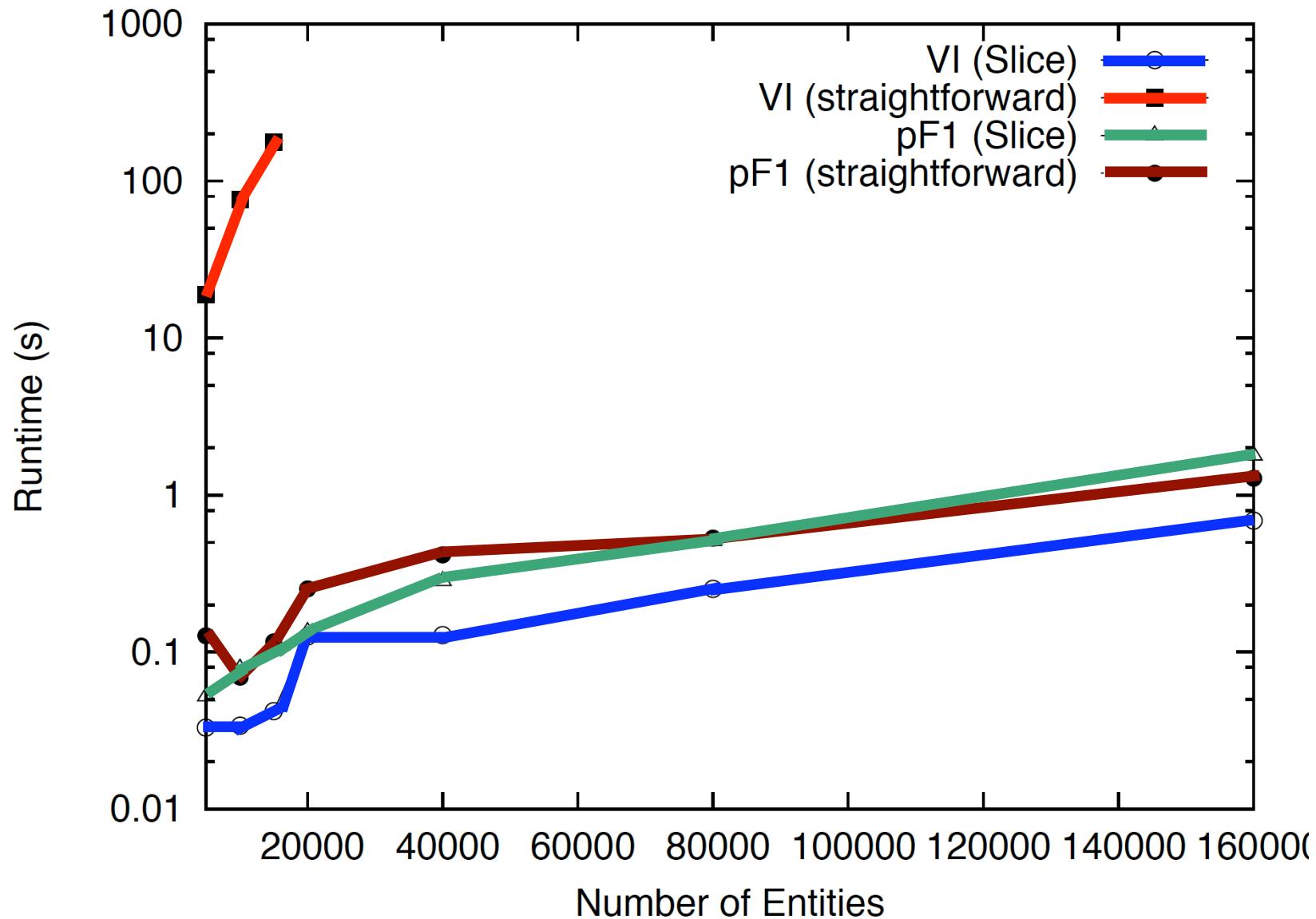
Road Map



Slice Runtime



Slice Runtime



Conclusion

- Existing measures conflict
- Generalized merge distance provides
 - Configurability to suit different applications
 - Framework for exploring relationships between measures
 - Efficient algorithm (Slice) for computing many distance measures

Thanks!

Relationships Between Measures

- Merge Distance: $\text{GMD}(R, G)$
where $f_m(x, y) = 1$, $f_s(x, y) = 1$
- Pairwise Recall: $1 - \text{GMD}(R, G)/\text{GMD}(\perp, G)$
where $f_m(x, y) = xy$, $f_s(x, y) = 0$
- Pairwise Precision: $1 - \text{GMD}(R, G)/\text{GMD}(R, \perp)$
where $f_m(x, y) = 0$, $f_s(x, y) = xy$
- Variation of Information: $\text{GMD}(R, G)$
where $f_m(x, y) = f_s(x, y) = h(x + y) - h(x) - h(y)$

$$h(x) = \frac{x}{N} \log \frac{x}{N}$$