Regular Expressions

Definitions

Equivalence to Finite Automata

RE's: Introduction

- *Regular expressions* are an algebraic way to describe languages.
- They describe exactly the regular languages.
- If E is a regular expression, then L(E) is the language it defines.
- We'll describe RE's and their languages recursively.

RE's: Definition

- Basis 1: If a is any symbol, then **a** is a RE, and L(**a**) = {a}.
  - Note: {a} is the language containing one string, and that string is of length 1.
- Basis 2: ϵ is a RE, and $L(\epsilon) = \{\epsilon\}$.
- Basis 3: $\emptyset$ is a RE, and $L(\emptyset) = \emptyset$.

RE's: Definition – (2)

- Induction 1: If E_1 and E_2 are regular expressions, then E_1+E_2 is a regular expression, and $L(E_1+E_2) = L(E_1) \cup L(E_2)$.
- Induction 2: If E_1 and E_2 are regular expressions, then E_1E_2 is a regular expression, and $L(E_1E_2) = L(E_1)L(E_2)$.

Concatenation : the set of strings wx such that w
Is in $L(E_1)$ and x is in $L(E_2)$.

RE's: Definition – (3)

- Induction 3: If E is a RE, then E* is a RE, and L(E*) = (L(E))*.

Closure, or "Kleene closure" = set of strings $w_1w_2...w_n$, for some $n \geq 0$, where each w_i is in L(E).

Note: when n=0, the string is ϵ.

Precedence of Operators

- Parentheses may be used wherever needed to influence the grouping of operators.
- Order of precedence is * (highest), then concatenation, then + (lowest).

Examples: RE's

- L(**01**) = {01}.
- L(**01**+**0**) = {01, 0}.
- L(**0**(**1**+**0**)) = {01, 00}.
  - Note order of precedence of operators.
- L(**0***) = {ε, 0, 00, 000,... }.
- L((**0**+**10**)*(ε+**1**)) = all strings of 0's and 1's without two consecutive 1's.

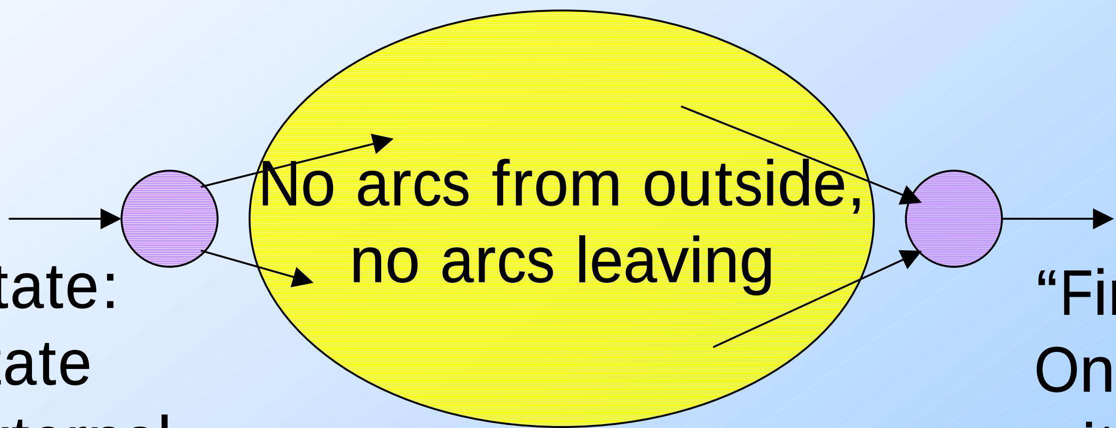
Equivalence of RE's and Automata

- We need to show that for every RE, there is an automaton that accepts the same language.
  - Pick the most powerful automaton type: the ϵ-NFA.
- And we need to show that for every automaton, there is a RE defining its language.
  - Pick the most restrictive type: the DFA.

Converting a RE to an ϵ-NFA

- Proof is an induction on the number of operators (+, concatenation, *) in the RE.
- We always construct an automaton of a special form (next slide).

Form of ϵ-NFA's Constructed

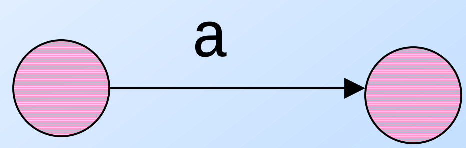


Start state:
Only state
with external
predecessors

"Final" state:
Only state
with external
successors

RE to ϵ-NFA: Basis

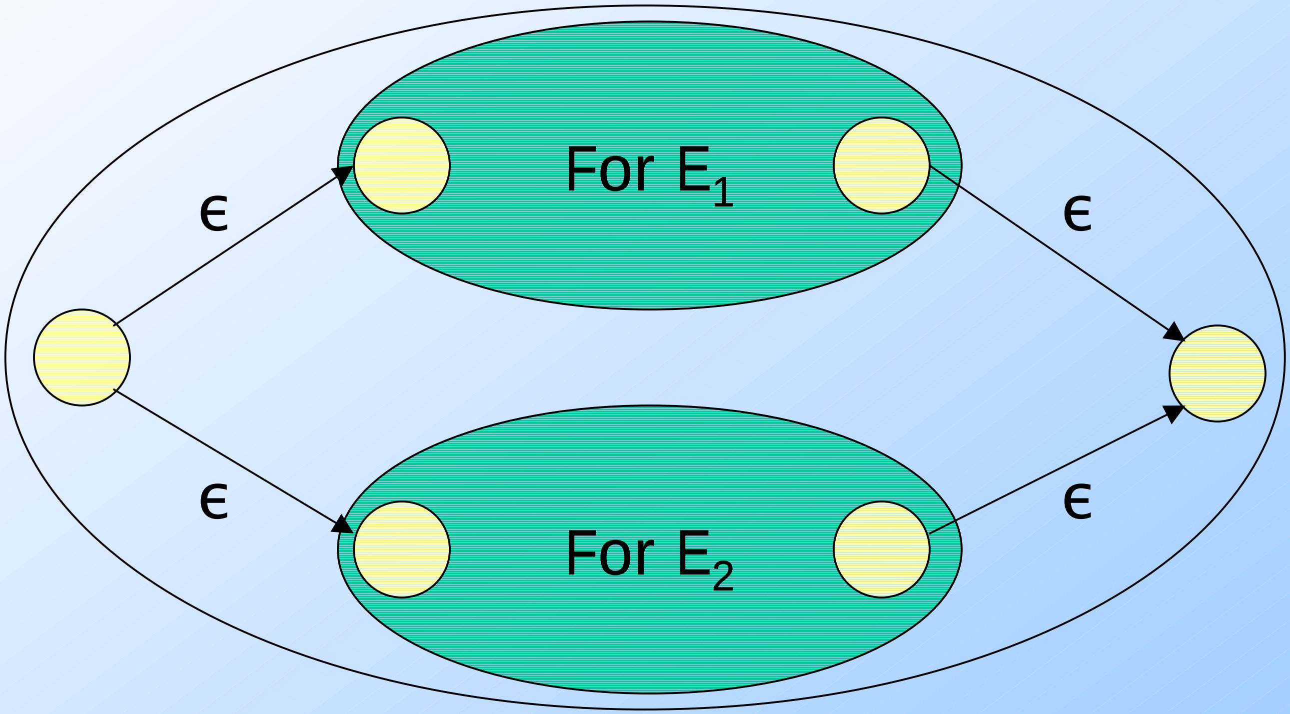
- Symbol **a**:



- ϵ:

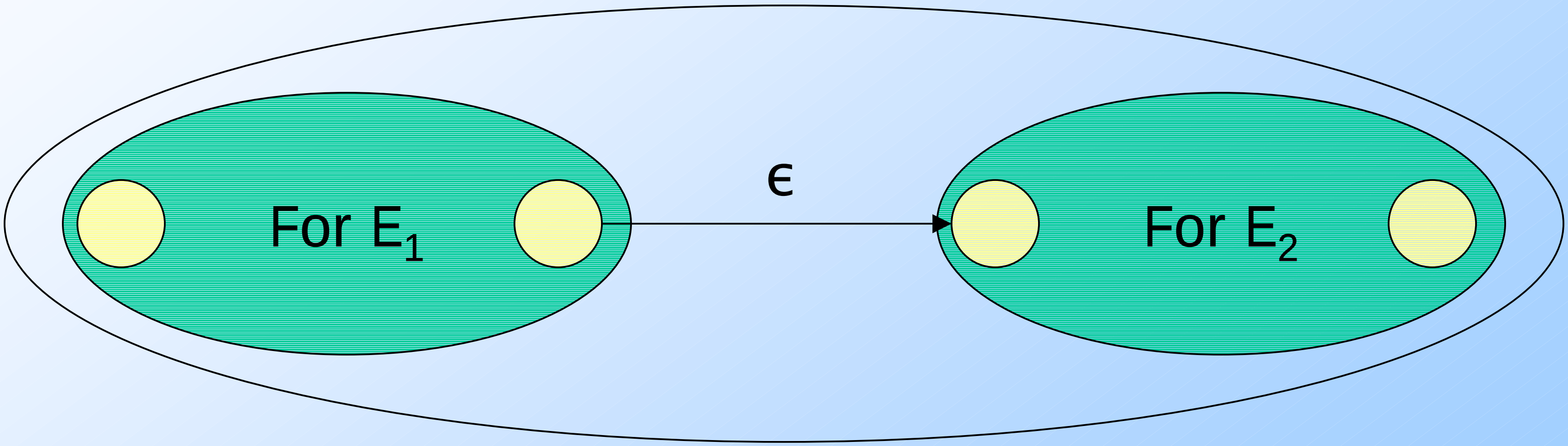
- ∅:

RE to ϵ-NFA: Induction 1 – Union



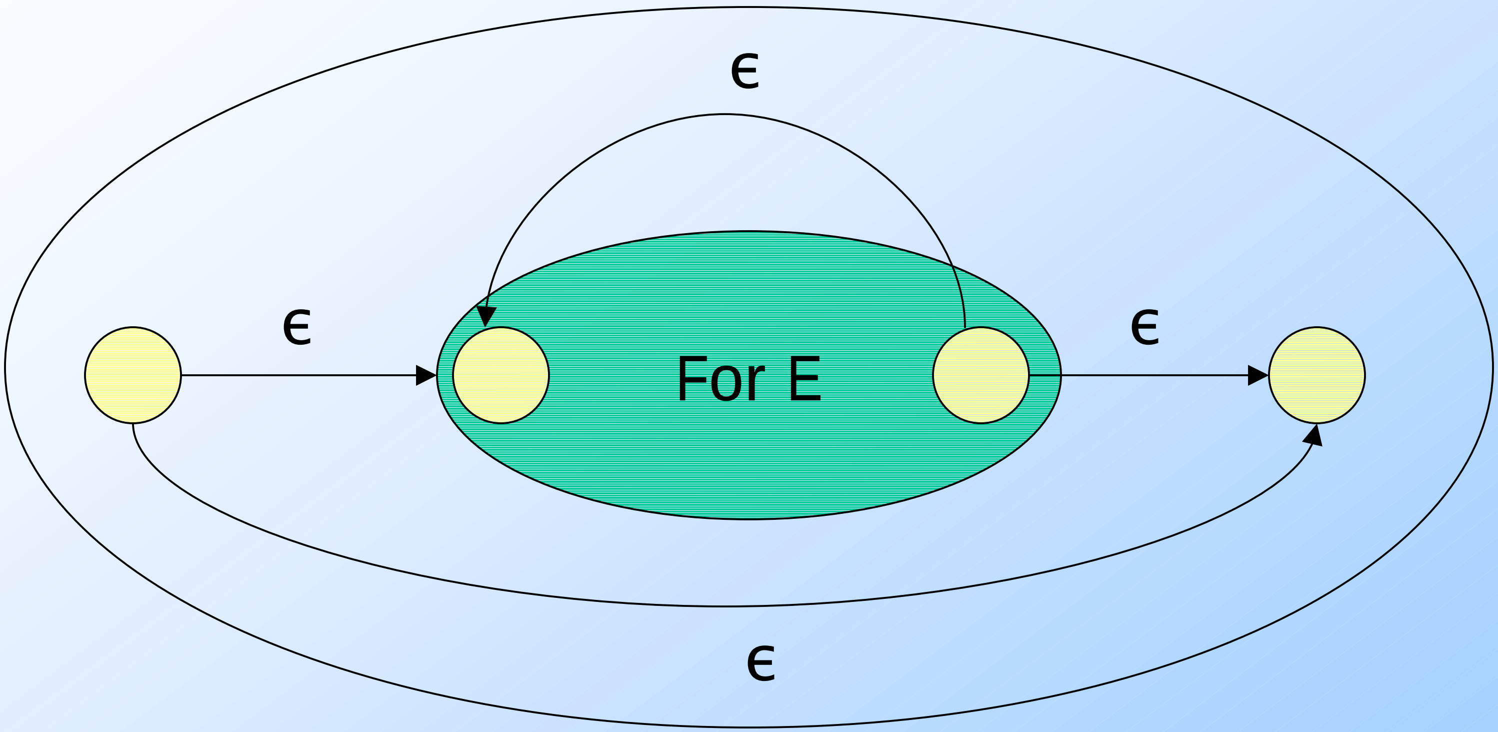
For $E_1 \cup E_2$

RE to ϵ-NFA: Induction 2 – Concatenation



For E_1E_2

RE to ϵ-NFA: Induction 3 – Closure



For E*

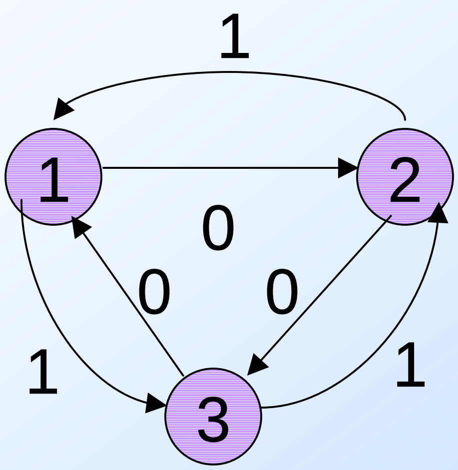
DFA-to-RE

- A strange sort of induction.
- States of the DFA are assumed to be 1,2,...,n.
- We construct RE's for the labels of restricted sets of paths.
  - Basis: single arcs or no arc at all.
  - Induction: paths that are allowed to traverse next state in order.

k-Paths

- A k-path is a path through the graph of the DFA that goes though no state numbered higher than k.
- Endpoints are not restricted; they can be any state.

Example: k-Paths



0-paths from 2 to 3:
RE for labels = **0**.

1-paths from 2 to 3:
RE for labels = **0**+**11**.

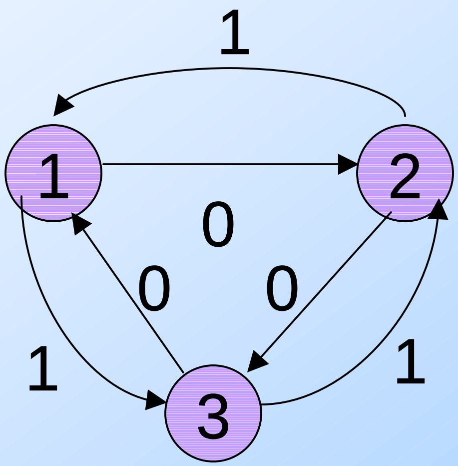
2-paths from 2 to 3:
RE for labels =
(**10**)***0**+**1**(**01**)***1**

3-paths from 2 to 3:
RE for labels = ??

k-Path Induction

- Let R_{ij}^k be the regular expression for the set of labels of k-paths from state i to state j.
- Basis: k=0. R_{ij}^0 = sum of labels of arc from i to j.
  - $\emptyset$ if no such arc.
  - But add ϵ if i=j.

Example: Basis



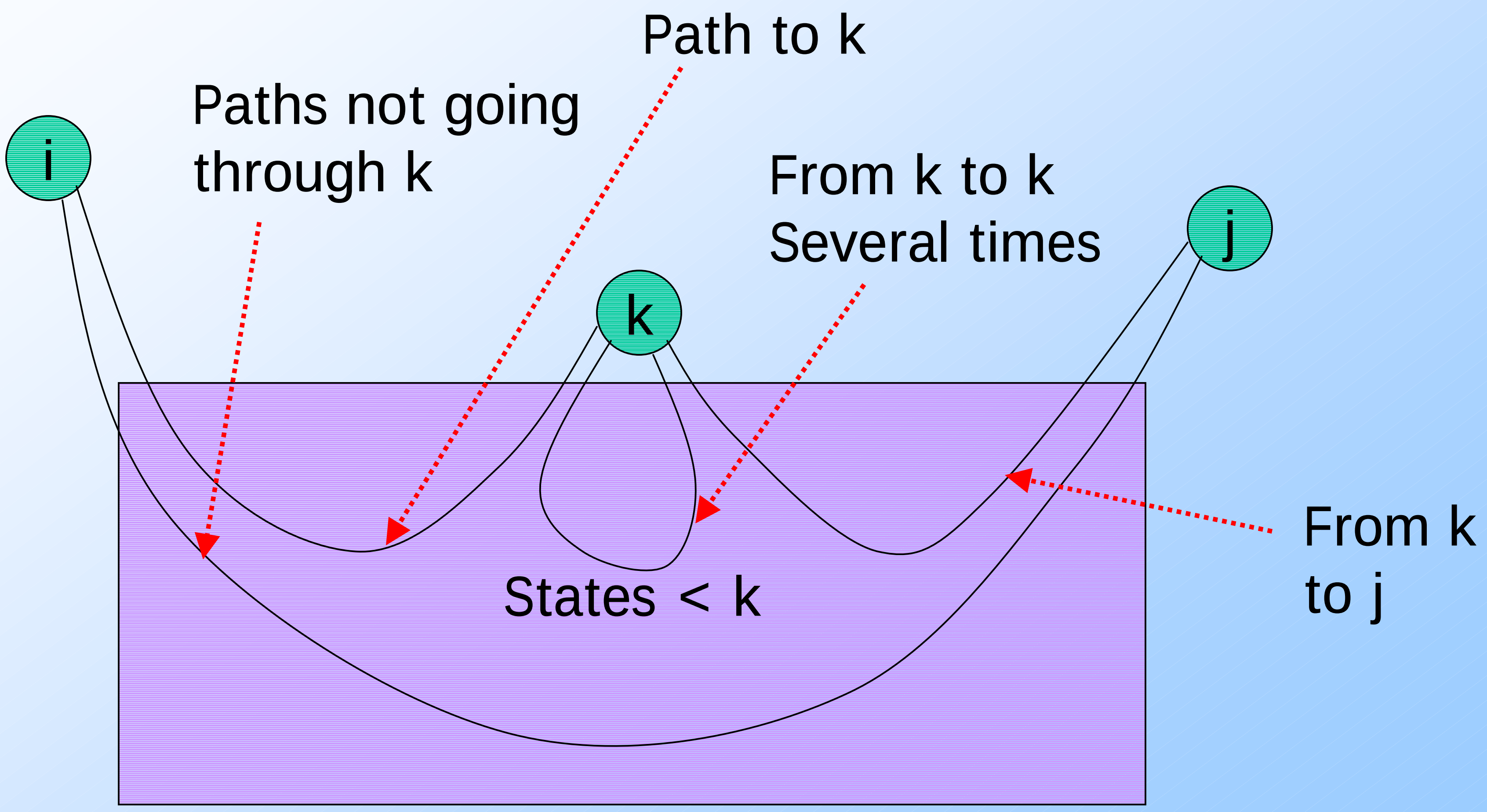
- $R_{12}{}^{0}$ = **0**.
- $R_{11}{}^{0} = \varnothing + \epsilon = \epsilon$.

k-Path Inductive Case

- A k-path from i to j either:
  1. Never goes through state k, or
  2. Goes through k one or more times.

$$R_{ij}^{k} = \underbrace{R_{ij}^{k-1}}_{\text{Doesn't go through k}} + \underbrace{R_{ik}^{k-1}}_{\text{Goes from i to k the first time}}\underbrace{(R_{kk}^{k-1})^*}_{\text{Zero or more times from k to k}}\ \underbrace{R_{kj}^{k-1}}_{\text{Then, from k to j}}.$$

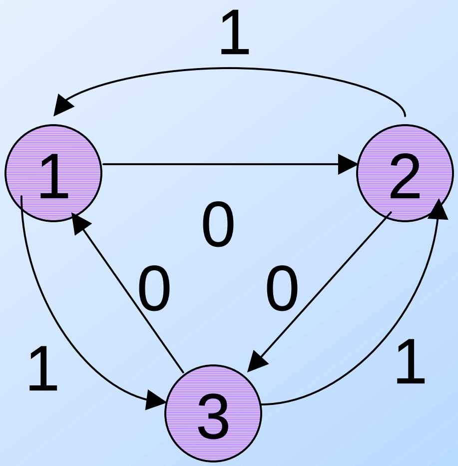
Illustration of Induction



Final Step

- The RE with the same language as the DFA is the sum (union) of R_{ij}^{n}, where:
  1. n is the number of states; i.e., paths are unconstrained.
  2. i is the start state.
  3. j is one of the final states.

Example



- $R_{23}^{3} = R_{23}^{2} + R_{23}^{2}(R_{33}^{2})^{*}R_{33}^{2} = R_{23}^{2}(R_{33}^{2})^{*}$
- $R_{23}^{2} = (\mathbf{10})^{*}\mathbf{0}+\mathbf{1}(\mathbf{01})^{*}\mathbf{1}$
- $R_{33}^{2} = \mathbf{0}(\mathbf{01})^{*}(\mathbf{1}+\mathbf{00}) + \mathbf{1}(\mathbf{10})^{*}(\mathbf{0}+\mathbf{11})$
- $R_{23}^{3} = [(\mathbf{10})^{*}\mathbf{0}+\mathbf{1}(\mathbf{01})^{*}\mathbf{1}][(\mathbf{0}(\mathbf{01})^{*}(\mathbf{1}+\mathbf{00}) + \mathbf{1}(\mathbf{10})^{*}(\mathbf{0}+\mathbf{11}))]^{*}$

Summary

- Each of the three types of automata (DFA, NFA, ϵ-NFA) we discussed, and regular expressions as well, define exactly the same set of languages: the regular languages.

Algebraic Laws for RE's

- Union and concatenation behave sort of like addition and multiplication.
  - + is commutative and associative; concatenation is associative.
  - Concatenation distributes over +.
  - Exception: Concatenation is not commutative.

Identities and Annihilators

- $\emptyset$ is the identity for +.
  - $R + \emptyset = R$.
- ϵ is the identity for concatenation.
  - $\epsilon R = R\epsilon = R$.
- $\emptyset$ is the annihilator for concatenation.
  - $\emptyset R = R\emptyset = \emptyset$.